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Specification and Drawings, as originally filed, with Application for Patent Serial No:
2,338,458, on February 27, 2001, by **IOAN DANCEA**, for "Method and VLSI Circuits Allowing
to Change Dynamically the Logical Behaviour".

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ABSTRACT

A method, named the product terms method that allows to implement and/or to change dynamically the logical behaviour of any combinational or synchronous sequential circuits has been presented. The method uses for every product term of logical equations, expressed as a sum-of-product, three memory words: mask word, product word and function word. The words of all product terms are ranged in a table, which characterize the logical behaviour of the circuit.

The invention provides the hardware structure of several new types of VSLI circuits, having re-configurable logic behaviours. A first embodiment implements any type of multiple output combinational circuit, a second embodiment implements any synchronous sequential circuit with only clock input and, a third embodiment implements any synchronous sequential circuits with data inputs and clock input.

An expert system capable to generate the tables used for the product terms method by interpreting and analysing the logical equations either supplied by the user or found in a database is also provided.

METHOD AND VLSI CIRCUITS ALLOWING TO CHANGE DYNAMICALLY THE LOGICAL BEHAVIOUR

BACKGROUND OF THE INVENTION

Field of invention

This invention is directed to logically modifiable VLSI digital circuits, and more specifically to combinational and synchronous sequential circuits that may be dynamically configured and re-configured to perform a digital function of choice.

Related art

Digital circuits are realised with integrated circuits. An integrated circuit, called chip, contains electrical components such as transistors, diodes, resistors and capacitors. The various components are interconnected inside the chip to form gates and memory elements.

A combinational circuit (CC) is obtained using exclusively elementary gates (NOT, OR, AND, XOR, NOR, etc). Each combinational circuit have a hardware structure designed to perform definite functions, as specified by logical equations. Consequently, there is a specific hardware structure for a 3 to 8 decoder, another hardware structure for a 4 bits comparator, and so on.

Sequential circuit, and the invention considers only the synchronous sequential circuits (SSC), contains gates and memory elements. Information is charged in memory elements to establish the initial state or the next state of the circuit. The model of synchronous sequential circuit in which the outputs depend on both the input variables and the state variables is called Mealy machine model. In another model of synchronous sequential circuit, called Moore machine model, the outputs depend only on the state variables. It is always possible to convert from one representation of synchronous sequential model to the other. For the purpose of the invention, any synchronous sequential circuit is represented as a Moore machine. Figure 1 illustrates the Moore machine model,

where combinational circuit **10** determines the next state stored in the memory part **25**, and combinational circuit **20** determines the outputs. In the general case, combinational circuit **10** is defined by equations including both, input variables and state variables. As it is known, an important number of usual synchronous sequential circuits define their next state by logical equations that use only the state variables, without considering any input variable, except the clock. For example, all the counters have this structure. The nature of the logical flip-flops used in the memory part **25** influences the structure of combinational circuit **10** and/or **20**. The simplest solution for the memory part **25** uses D synchronous flip-flops. Like combinational circuits, each synchronous sequential circuit has its own hardware structure. Thus, there is a specific hardware structure for an 8 bits synchronous binary counter, another hardware structure for a 6 bits pseudo-random generator, and so on.

As the digital technique advanced, the design of circuits was realised by developing VLSI (very large scale integrated) technologies. A general digital VLSI circuit consists of logic functional blocks and memory elements, which are interconnected by connection lines to perform a specific task (for example a VLSI processor). Many real applications need to configure a specific logical structure, or to be able to change dynamically an initial structure, consequently the ability to cover a re-configurable logical behaviour. These kinds of applications use programmable VLSI logic devices. Programmable logic devices have progressed from simple group of gates to very complex Field Programmable Gate Arrays (FPGAs), which have a large number of inputs and outputs, memory elements, programmable logic blocks, tracks for interconnecting the logical blocks to each other and to inputs and outputs. The FPGA's are used to implement or to modify structures with a high number of logic functions, which result in lower part counts, lower power dissipation, higher speed of operation and greater flexibility than if discrete components are used.

Using the conventional programmable VLSI circuits, to realise a particular logical structure or to modify such a structure and implicitly its logical behaviour, requires, for any combinational or synchronous sequential complex circuit, a

programming effort and a dedicated programmable device. Consequently, a separate program is needed each time such a particular circuit is implemented, or modified in hardware. In addition, for complex circuits the developer must add the logic part for testing the functional structure of the circuit, which takes a large surface of the real estate.

SUMMARY OF THE INVENTION

The present invention seeks to overcome the disadvantages of the prior art associated with digital programmable VLSI circuits.

It is an object of the present invention to provide a method, named the product terms method, which allows the implementation, by software and/or by hardware, of multiple output combinational circuits and synchronous sequential circuits.

It is another object of the present invention to provide the hardware structure of several new types of VLSI circuits, having re-configurable logic behaviours. A first embodiment of the invention can implement any type of multiple output combinational circuit, a second embodiment can implement any synchronous sequential circuit with only clock input and, a third embodiment can implement any synchronous sequential circuits with data inputs and clock input, known as Moore machine model.

It is another object of the present invention to provide an expert system capable to generate the tables used for the product term method by interpreting and analysing the logical equations either supplied by the user or found in a database.

According to one aspect of the invention a VLSI circuit for implementing in hardware any multiple output combinational circuit is provided, having m inputs, r outputs and the behaviour expressed in logical sum-of-product equations with n product terms $p(k)$. This VLSI circuit comprises a register with m bits for storing the input variables, n basic cells, where a cell $C(k)$ is used for determining the logical value of a product term $p(k)$ of the equations for given inputs, and a logical

summing circuit, realised with r OR gates, each one with n inputs, associated with the cell $C(k)$ for receiving the logical value of product terms $p(k)$ and outputting the r bits of output functions.

According to another aspect of the invention a VLSI circuit for implementing in hardware any synchronous sequential circuit with clock input only and outputs taken from the state register is provided, having the next state functions expressed in logical sum-of-product equations with n product terms $p(k)$. This VLSI circuit comprises a clock input, a state register with m bits for storing the state variables, n basic cells, where a cell $C(k)$ is used for determining the logical value of a product term $p(k)$ of the equations for the next state, a logical summing circuit, realised with m OR gates, each one with n inputs, associated with the cell $C(k)$ for receiving the logical value of product terms $p(k)$ and outputting the m bits of the next state functions, and a feedback connection to establish the next state.

According to a further aspect of the invention a VLSI circuit for implementing in hardware any synchronous sequential circuit with data inputs and clock input is provided. It represents a special association of the VLSI circuits mentioned above, precisely the VLSI circuit for implementing in hardware any multiple output combinational circuit and the VLSI circuit for implementing in hardware any synchronous sequential circuit with clock input only.

According to still another aspect of the invention an expert system for analysing and interpreting the logical behaviour of digital circuits, given by input variables, outputs and logical equations in form of sum-of-products is provided, where analysing and interpreting means to validate the input variables and the outputs, to analyse the syntactic correctness of logical equations, to generate for every product term the mask word, the product word and the function word and to transform the individual lists, generated for each product term in a single memory list defining the logical behaviour of a digital circuit and storing it in a database.

According to still another aspect of the invention a method for dynamically configuring the logical behaviour of a VLSI circuit is provided by performing the

steps of inputting the input variables, the outputs and a plurality of sum-of – product equations, which describe the logical behaviour of a digital circuit, generating three memory words uniquely defining each product term of the equations and storing the mask word, the product word, and the function word into a corresponding register associated with a cell $C(k)$.

Advantageously, using a fixed hardware structure, the invention provides a mechanism for implementing a large number of VLSI digital circuits having each one a specified logical behaviour. This logical behaviour can be dynamically changed, without replacing the physical structure and/or without using any programming task. One basic cell with a particular physical structure is provided for implementing one product term of any combinational or synchronous sequential digital circuit. For changing the functions performed by the circuit, a set of weights uniquely defining the desired circuitry functions are calculated with high accuracy by an expert system and loaded in the registers of the basic cells. No testing of the logical structure realised with the VLSI circuit is necessary, and this significantly reduces manufacturing costs.

The "Summary of the Invention" does not necessarily disclose all the inventive features. The invention may reside in a sub-combination of the disclosed features.

BRIEF DESCRIPTION OF THE DRAWINGS

The invention will be now explained by way of example only and with reference to the following drawings, wherein similar references are used in different figures to denote similar components:

Figure 1 shows the Moore machine model for synchronous sequential circuits (prior art);

Figure 2A is a first example of how the mask word, product word and function word are applied to the input variables to provide the outputs, according to the invention;

Figure 2B shows another example, for illustrating the method of the invention;

Figure 3 shows the structure of a VLSI circuit for implementing a plurality of combinational circuits;

Figure 4 shows the structure of a 8 bits basic cell;

Figure 5 is a block diagram of the system used for configuring the VLSI circuit of Figure 3;

Figure 6 shows the structure of a VLSI circuit for a plurality of synchronous sequential circuits with clock input and outputs from the state register;

Figure 7 shows the structure of a VLSI circuit for general Moore machine model, implementing a plurality of synchronous sequential circuits with data inputs and clock input;

Figure 8 shows the association of two VLSI circuits to form a 4 bits fast adder;

Figure 9 shows the association of three VLSI circuits to form a structure, which can implement several types of direct and reverse counters, with outputs displayed over a 7-segment indicator;

Figure 10 shows the flowchart of the expert system that automatically generates the values of weights used in the product terms method.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

The following description is of preferred embodiments by way of example only and without limitation to combination of features necessary for carrying the invention into effect.

The present invention proposes a novel solution for implementing a combinational and/or synchronous sequential digital circuit, having a one and only one physical structure and able to implement a variety of sum-of-product functions. The method according to the invention is called *the product terms method*. More precisely it implements in hardware the general pattern of *the*

sum-of-product form of a set of logical equations defining a multiple output combinational circuit, or the general pattern of *the sum-of-product form of several sets of logical equations* defining a synchronous sequential circuit. This is contrasting with the prior art solution, which uses direct implementation of equations by gates, to designate each logical circuit. To minimize the number of required gates of the circuit, it is recommended that the number of product terms in logical equations describing the behaviour of any multiple output combinational circuit and/or synchronous sequential circuit, be kept to a minimum. The minimization process is outside the scope of the present invention.

Let's consider the logical equations in the form of sum-of-products, that describes the behaviour of a combinational circuit. First, the input variables and the outputs, are placed in succession on virtual registers, and this placement attaches to each of them a weight that has a value expressed as a power of 2 starting with 2^0 (right alignment). Second, analysing successively all product terms, a table is generated. To construct this table, each product term provides information through three memory words. The first, called the *mask word*, corresponds to the sum of the weights of input variables representing a product term. Next, the *product word* represents the sum of the weights of the variables of a product term that give a value of logic 1 to the analysed product term. The third, called the *function word*, represents the sum of the output weights where a product term has a value of logic 1.

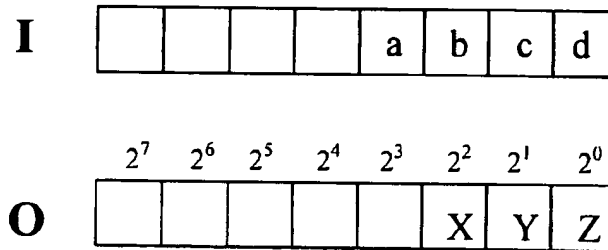
Using this numerical weights, it can be realised an indirect manipulation of the input variables, product terms and outputs. Because each product term may have a logical value of 1 or 0, in order to find this value for a set of given input variables, the method uses the following steps: an AND masking with the mask word to extract the variables of a product term, followed by a comparison operation with the product word to determine its real value, 1 or 0. Finally, because a product term may be present in several equations, the function word of an active product term (its logical value is 1) is OR-ed with function words of other actives product terms to determine the outputs. Since the general sum-of-products form of logical equations is similar for a large diversity of circuits, one

fixed hardware structure can be used for all circuits having this similar form of equations.

The first example relates to a multiple output combinational circuit, where the product terms method is used to obtain the outputs associated with the input variables, and with the words defining the logical equations. It is to be understood that the values of the words are taken by way of example, for explaining how the functions are represented. In this first example the words have 8 bits (a byte). It is to be understood that the size of the word is selected according to the application.

The input variables (also called inputs) are denoted with a , b , c and d , and the output functions (also called outputs), with X , Y , and Z . The inputs are right aligned and "packaged" into an input word "I", while the outputs are right aligned and "packaged" into an output word "O", each word having 8 bits (a byte). The non-occupied places in each word are filled with the bit 0 and, for simplicity, their contribution to the sum of weights (always 0), will be ignored.

This can be shown as below:



Let's assume that for this first example a multiple output combinational circuit is characterized by the following sum-of-products equations, which give the outputs X , Y , and Z as a combination of inputs a , b , c and d :

$$X = a \cdot b + c \cdot d ,$$

$$Y = a \cdot b + \bar{c} \cdot d , \text{ and}$$

$$Z = \bar{a} \cdot \bar{b} \cdot \bar{c} + a \cdot b + \bar{c} \cdot \bar{d}$$

(1)

The functions comprise in this example five different product terms, namely $a \cdot b$; $c \cdot d$; $\bar{c} \cdot d$; $\bar{a} \cdot \bar{b} \cdot \bar{c}$; and $\bar{c} \cdot \bar{d}$. Each mask word is determined by the sum of the weights of input variables that compose a product term. The mask words are therefore, for the above example:

Mask word for the product term $a \cdot b = 2^3 + 2^2 = 12$

Mask word for the product term $c \cdot d = 2^1 + 2^0 = 3$

Mask word for the product term $\bar{c} \cdot d = 2^1 + 2^0 = 3$ (2)

Mask word for the product term $\bar{a} \cdot \bar{b} \cdot \bar{c} = 2^3 + 2^1 + 2^0 = 14$

Mask word for the product term $\bar{c} \cdot \bar{d} = 2^1 + 2^0 = 3$

Each product word represents the sum of the weights of the input variables that give a logical value 1 for the analysed product term. The product words are therefore, for the above example:

Product word for the product term $a \cdot b = 2^3 + 2^2 = 12$

Product word for the product term $c \cdot d = 2^1 + 2^0 = 3$

Product word for the product term $\bar{c} \cdot d = 2^0 = 1$ (3)

Product word for the product term $\bar{a} \cdot \bar{b} \cdot \bar{c} = 0$

Product word for the product term $\bar{c} \cdot \bar{d} = 0$

The function word indicates the contribution of the product term into the outputs. For example, $a \cdot b$ is a product term that can give the logical value 1 to the output X, which has range 2^2 in the output word. The same product term $a \cdot b$ can give the logical value 1 to output Y, which has range 2^1 in the output word and can give the logical value 1 to output Z, which has range 2^0 . The function words are therefore, for the above example:

Function word for the product term $a \cdot b = 2^2 + 2^1 + 2^0 = 7$

Function word for the product term $c \cdot d = 2^2 = 4$

Function word for the product term $\bar{c} \cdot d = 2^1 = 2$ (4)

Function word for the product term $\bar{a} \cdot \bar{b} \cdot \bar{c} = 2^0 = 1$

Function word for the product term $\bar{c} \cdot \bar{d} = 2^0 = 1$

The individual values of all these words are summarized in Table 1.

Product Term	Mask word	Product word	Function word
$a \cdot b$	12	12	7
$c \cdot d$	3	3	4
$\bar{c} \cdot d$	3	1	2
$\bar{a} \cdot \bar{b} \cdot \bar{c}$	14	0	1
$\bar{c} \cdot \bar{d}$	3	0	1

Table 1.

Finally, each product term is associated with the respective mask, product and function values to form a single memory list, as in Table 2 below.

$a \cdot b$	$c \cdot d$	$\bar{c} \cdot d$	$\bar{a} \cdot \bar{b} \cdot \bar{c}$	$\bar{c} \cdot \bar{d}$
<u>12, 12, 7,</u>	<u>3, 3, 4,</u>	<u>3, 1, 2,</u>	<u>14, 0, 1,</u>	<u>3, 0, 1</u>

Table 2.

Figure 2A shows how the mask words, product words and function words are applied to the associated inputs to provide the outputs for a first numerical example. The bi-directional arrows indicate logical operations between the bits, while a unidirectional arrow shows the result of logical operations. In this

example, the logical functions (1) and the input byte "IN = 01h" are considered (h means hexadecimal, therefore IN=0000 0001 in binary, namely $a = 0$, $b = 0$, $c = 0$, and $d = 1$). The output value in this case is "OUT=03h", because:

$$X = 0 + 0 = 0$$

$$Y = 0 + 1 = 1$$

(5)

$$Z = 1 + 0 + 0 = 1$$

The columns Col.1 to Col 5 show the operations over each product term $a \cdot b$, $c \cdot d$, $\bar{c} \cdot d$, $\bar{a} \cdot \bar{b} \cdot \bar{c}$ and $\bar{c} \cdot \bar{d}$.

The first row, Row 1, shows the inputs and the second row, Row 2, shows the mask words as given by Table 1. The first intermediate results are shown in Row 3. These intermediate results are obtained by bit by bit AND operation (logical multiplication) between the inputs and the respective mask word.

Row 4 shows the product words associated with each product term, as given in Table 1. The second intermediate results are shown in Row 5. These intermediate results are obtained by EQUIVALENCE operations, which gives logic 1 when both bits are identical (i.e. both are logic 0 or logic 1) and gives a logical 0 when the bits are different, i.e. one is logic 0 and the other is logic 1 (EQUIVALENCE and $\overline{\text{XOR}}$ express the same logical operation).

Row 6 show AND operations between all the bits find in a colon at the place of second intermediate results (Row 5), to determine if for the considered values of inputs, the considered product term has a logic value of 1 or 0. All the functions words of Row 7, for which the product term has the logic value of 1, therefore the Row 6 indicate the logical 1, are OR-ed to obtain the final result at outputs.

When maintaining the same number of inputs and same number of outputs, any change in the logical equations requires only a change of the values in the table without changing the physical structure of the VLSI circuit. The new tables provide the logical behaviour for the new circuit.

Suppose, that the name and the number of input variables, e.g. a, b, c, d , and outputs, e.g. X, Y, Z , are not changed and the following new equations are considered:

$$\begin{aligned} X &= a \cdot b + c \cdot d, \\ Y &= a \cdot b + c \cdot \bar{d}, \text{ and} \\ Z &= \bar{a} \cdot \bar{b} \cdot c + a \cdot b + \bar{c} \cdot d, \end{aligned} \quad (6)$$

Table 3, reflecting the circuit behaviour, is constructed as Table 2 above.

$a \cdot b$	$c \cdot d$	$c \cdot \bar{d}$	$\bar{a} \cdot \bar{b} \cdot c$	$\bar{c} \cdot d$
<u>12, 12, 7,</u>	<u>3, 3, 4,</u>	<u>3, 2, 2,</u>	<u>14, 2, 1,</u>	<u>3, 1, 1</u>

Table 3.

For this second example, when the input "IN=01h" the output is, if we calculate the function (6):

$$\begin{aligned} X &= 0 + 0 = 0 \\ Y &= 0 + 0 = 0 \\ Z &= 0 + 0 + 1 = 1 \end{aligned} \quad (7)$$

Figure 2B shows how the mask words, product words and function words are applied to the associated inputs to provide the outputs for a second numerical example, according to equations (6).

Thus, the advantage of the product terms method is the use of one and only one hardware structure to simulate/implement a large number of combinational circuits, in which each circuit has its own table specifying its logical behaviour.

The invention will be described in connection with Figures 3 and 4. Figure 3 shows the structure of a VLSI circuit 40 for implementing a plurality of multiple output combinational circuits, according to the invention. Figure 4 illustrates an

eight bit basic cell **30**. A basic cell, as in Figure 4, is needed for each product term appearing in the logical equations. The succession of the operations executed by the VLSI circuit **40**, of Figure 3, including the basic cell of Figure 4, will be explained in conjunction with the numerical example presented in Figure 2A, for the product term "*a.b*" shown in the first column (Col 1). For the general case, we denote by "*k*" the index of a cell **30-k**, where "*k*" takes values between 1 and "*n*".

As discussed in connection with Figure 2A, the mask words, product words and function words for a certain multiple output combinational circuit (here the logical equations given by (1)) are predetermined in software using an expert system as it will be later explained. The words, which characterize a product term, in this example *a.b*, are placed in registers **32**, **33** and **34**, respectively. As given by Table 1, the mask word for Col. 1 is 12 (binary 0000 1100); the product word is 12 (binary 0000 1100); and the function word is 7 (binary 0000 0111). In the same time, the input variables (binary 0000 0001 in this example) are placed in the IN (input) register **31**.

The input variables are AND-ed with the mask word using AND gates **35**. Because in this example the words have 8 bits, eight 2 inputs AND gates are needed, each one for implementing the AND operation between a bit of the input with a corresponding (in range) bit of the mask word.

The first intermediate result of this operation is shown in Col. 1, Row 3, of Figure 2A, and it is 0000 0000. Next, the logical value of the product term (1 or 0) is determined by a comparison operation (bit by bit) between the product word 0000 1100 and the intermediate result obtained above. As indicated, this comparison operation uses EQUIVALENCE gates **36**, marked with $\overline{\text{XOR}}$ on Figure 4. Each EQUIVALENCE gate has two inputs, one for a bit of the product word and one for the corresponding bit of the first intermediate result. The outputs of the EQUIVALENCE gates, i.e. the second intermediate result is 111 0011.

The logical value of the product term is then determined by AND-ing the 8 bits of the comparison result using the unique 8 inputs AND gate **37**. If the

output of this gate is at logical 1, it means that for the input variables considered, the corresponding product term is active and thus it has the logical value 1 in all equations where it is present. In Col 1 of the considered example, the output of the AND gate 37, for the product term $a.b$ and the input $IN = 01h$, has the logical value 0, and therefore the contribution of this product term is 0 to all equations where it is present. If the output of AND gate 37 is 1, like that corresponding to the product terms found in Col 3 and in Col 4 in the considered example, this signal opens the eight 2 inputs AND gates 38, permitting to the corresponding function words to be OR-ed in the OR gates 39, to determine the circuit outputs. There is only one group of OR gates for the VLSI circuit 40 of Figure 3, each gate having "n" inputs, where "n" expresses its maximum capacity in number of basic cells. If, for a specified VLSI circuit 40, "n" becomes very large, the final OR gates may be realised in a multi-level solution.

The gates in all cells 30-k act as a filter between the information that characterizes the product terms and the outputs, for implementing in parallel the product terms method. It must be noted that the input register 31 and the OR gates 39 are not part of the basic cell 30.

At the initialisation level the $(3n)$ registers in the "n" basic cells are charged with the information that corresponds to "n" product terms, from a table like Table 2 generated by an expert system according to selected logical equations. When changing that information, the logical behaviour of the VLSI circuit 40 changes accordingly, but the physical structure of the VLSI circuit 40 remains unchanged.

To successfully implement a plurality of functions on a VLSI circuit 40, and to allow flexibility in changing the functions as needed, the total number "n" of basic cells 30 must be increased as much as permitted by the available technology.

For example, a VLSI circuit 40 having 8 inputs, 8 outputs and $n = 100$ basic cells, therefore $(3 \times 100 + 1)$ registers, can implement any combinational circuit with no more than 8 inputs, and with no more than 8 outputs and with no more than "n" product terms. A non-exhaustive list of these types of

combinational circuits is presented below. The number of product terms was approximated for some of the circuits.

Two 2 to 4 binary decoders (4 inputs, 8 outputs, 8 product terms).

One 3 to 8 binary decoder (3 inputs, 8 outputs, 8 product terms).

One 8 to 3 binary priority encoder (8 inputs, 3 outputs, 9 product terms).

One 4 bits comparator (8 inputs, 3 outputs, 46 product terms).

One 4 to 1 multiplexer (6 inputs, 1 outputs, 4 product terms).

Two 2 to 4 demultiplexers (6 inputs, 8 outputs, 8 product terms).

One 3 to 8 demultiplexer (4 inputs, 8 outputs, 8 product terms).

One 4 bits fast binary parallel adder (8 inputs, 5 outputs, approximately 80 product terms).

One 2 digits BCD parallel adder (8 inputs, 8 outputs, approximately 80 product terms).

One 4 bits integer multiplier (8 inputs, 8 outputs, approximately 90 product terms).

One 2 digits BCD multiplier (8 inputs, 8 outputs, approximately 90 product terms).

One 1 digit code converter BCD to 7 segments (4 inputs, 7 outputs, 16 product terms).

As mentioned before, all registers **32**, **33** and **34** that are not used for a certain application are loaded with logic 0 at the initialisation level. Evidently, a VLSI circuit **40** having $(3n + 1)$ registers, cannot implement functions with more than "n" product terms.

The VLSI circuit **40** can be configured in an environment, which has a very simple processor, or a typical personal computer (PC) such as PCs manufactured by IBM Company, Apple Company, or any other manufacturer. A block diagram of such a system is shown in Figure 5. The user interrogates an expert system **80**, which gives him the possibility either to introduce the logical equations **79** or to consult a database **77**, where these equations are stored. The expert system generate one or several tables containing words (mask words (MW), product words (PW) and function words (FW)), which characterize the

circuit behaviour, and which are charged dynamically in the registers of type **32**, **33** and **34** belonging to the VLSI circuit **40**. Next, the user introduces or modifies the inputs **78**. Any time when the user modifies the logical equations **79** he modifies the logical behaviour of the circuit, and any time when he modifies the inputs **78** the VLSI circuit **40** produces new outputs for the selected circuit. A control program can replace the human user and, therefore, the process of dynamically changing the logical behaviour of any circuit can be performed automatically. There are two major solutions to realise physically the VLSI circuit **40**. First, the registers of the VLSI circuit **40** are of type random access memory (RAM) and in this case the VLSI circuit **40** must be permanently tie up to computer to have the possibility of frequently changing its logical behaviour. Second, the registers of the VLSI circuit **40** are of type programmable read only memory (PROM) or re-programmable read only memory (REPROM) and in this case, previously a logical behaviour established, the VLSI circuit **40** can be removed (like a FPGA programmable logic devices) and can be used to implement a specific task. Eventually, it is reprogrammed later, when its logical behaviour must be changed. An arrangement as in Figure 5, with minor modifications, can also be used for sequential implementations of the product terms method of the invention, for multiple output combinational or synchronous sequential circuits, using software.

Figure 6 shows the structure of a VLSI circuit **50** implementing a plurality of synchronous sequential circuits with clock input only, according to the invention. The VLSI circuit **50** uses the same eight bits basic cell **30** as illustrated in Figure 4. As mentioned earlier, it will be understood that the invention is not limited to this values for the memory part. This VLSI circuit **50** is very similar to the circuit of Figure 3, using the state register **201** instead of the input register **31**. As shown in Figure 6 the VLSI circuit **50** needs a clock (clk) input. A feedback **202** generates the next state at every clock input. The outputs of the circuit are taken directly from the outputs of the state register **201** and thus this VLSI circuit **50** has only one combinational circuit, used to determine the next state. Considering these modifications, the synchronous sequential VLSI circuit **50** of

Figure 6 works in a similar manner than that described for the combinational VLSI circuit **40** of Figure 3.

A VLSI circuit **50** like that presented in Figure 6, supposed to have $n = 20$ basic cells, thus $(3 \cdot 20 + 1)$ registers, can implement any synchronous sequential circuit with only clock input having no more than 8 bits state, and less than 20 product terms in logical equations. A non exhaustive list of these type of synchronous sequential circuits are (the exact number of product terms has not been determined for all circuits):

Any four 2 bits, or two 3 bits, or two 4 bits, or one 5 bits, or one 6 bits, or one 7 bits or one 8 bits direct binary counter (13 product terms for 8 bits counter).

Any four 2 bits, or two 3 bits, or two 4 bits, or one 5 bits, or one 6 bits, or one 7 bits or one 8 bits reverse binary counter (13 product terms for 8 bits counter

A two digit BCD direct counter (9 product terms).

A two digit BCD reverse counter (9 product terms).

Any modulo 3, or modulo 5, or modulo 6, or modulo 7, or modulo 9, or modulo 11, or modulo 12, or modulo 13, or modulo 14, or modulo 15 counter.

Any 2 bits, or 3 bits, or 4 bits counter, counting in any numeration system (for example a 4 bits counter which counts in excess 3 code).

Any shift left register with maximum 8 bits.

Any shift right register with maximum 8 bits.

Any left recirculating register with maximum 8 bits.

Any right recirculating register with maximum 8 bits.

Any 3 bits, or 4 bits, or 5 bits, or 6 bits, or 7 bits binary pseudo-random generator.

A minute chronometer (18 product terms).

A dice game (14 product terms).

Figure 7 illustrates an example of a general Moore synchronous sequential circuit **65** with two combinational circuits, precisely the Next state circuit **60** and the Combinational circuit' **70** for outputs. Only one cell **30** and one

cell **30'** are illustrated on each of the combinational circuits **60** and **70** for simplification. It is to be understood that, as in the previous examples, the circuit **60** may have a number of "n" basic cells **30** and the circuit **70** may have a number of "n'" basic cells **30'** for implementing a structure with "n" product terms in the next state equations, and "n'" product terms in the output equations. Unlike the previous examples, the basic cells **30** of the next state equations operate on 16 bits, but the basic cells **30'** of output equations operate on 8 bits. In this example, the input register **31** has a capacity of 8 bits and the state register **51**, has also a capacity of 8 bits.

As shown in Figure 7 the Next state circuit **60** needs a clock (clk) input. The state register **51** receives next state values from the OR-gate **39** on line **59**.

The VLSI circuits **40**, **50** and **65** presented in these preferred embodiments can be combined to allow passing the information in a way similar to that known inside a classical network composed by several level of gates. When the necessary number of product terms is too large for a given VLSI circuit of type **40**, **50** or **65**, two or several VLSI circuits of type **40**, **50** or **65** may be used and connected in series.

For example, the VLSI circuit **40** of Figure 3 (considered with 100 basic cells) used to implement a fast 4 bits binary parallel adder, can be replaced by two VLSI circuits of type **40**, precisely **40-1** and **40-2** as shown in Figure 8. Each VLSI circuit of type **40** of Figure 8 is similar to the VLSI circuit **40** of Figure 3. Circuits **40-1** and **40-2** of Figure 8 has only $n = 30$ basic cells (every 2 bits fast binary adder needs in logical equations 27 product terms). The VLSI circuit **40-1** performs the addition of first two bits, A_0 with B_0 and A_1 with B_1 , and outputs on lines **43**, **44** the bits of sum S_0 and S_1 and on line **45** the bit of carry C_2 . The VLSI circuit **40-2** performs the addition of last two bits, A_2 with B_2 and A_3 with B_3 , and outputs on lines **46**, **47** the bits of sum S_2 and S_3 and on line **48** the bit of carry C_4 .

A second example is presented in Figure 9, where it was associated a VLSI circuit **50** with two VLSI circuits of type **40**, precisely **40-3** and **40-4**, to obtain a structure which can implement direct and reverse counters with outputs

that use 7 segments display. For this last example, if all three VLSI circuits are implemented in a single chip, the two input registers **31** of the structure specified in Figure 3 for the VLSI circuits **40**, can be omitted because the state register **201** of the VLSI circuit **50** can be used as the input register for the VLSI circuit of type **40**.

Finally, we consider that a large number of connected VLSI circuits of the types presented in this invention, can directly execute programs at higher speeds, creating, in this way, a special data flow computer.

Figure 10 illustrates the flowchart of an expert system **80** capable to generate the tables used in the product terms method by interpreting and analysing the logical equations supplied by the user or found in a database.

When a request for a specific circuit structure arrives in step **81**, the expert system interacts in step **82** with the user and offers him a choice, either:

- inserting the behaviour of the selected circuit in the form of input variables, outputs and logical equations of the sum-of-products by the step **87**, or
- consulting a database storing circuit behaviours associated with a large number of configurations of physical structures, and retrieving the equations associated with the desired circuit.

By the step **83** the user consults a knowledge base by means of a classical *backward-chaining inference engine*. The inference engine works in a query mode, matching the user's answer to determine the nature of the required structure. The knowledge base contains the parametric description of each logical structure having a file in database.

Once the structure is identified, the system loads from the database **84** the appropriate file describing the circuit behaviour that is the input variables, the outputs, and the equations. Usually, such a file contains the description of several circuits intended to characterize a complex structure.

The most important step **85** of the expert system generates for any multiple output combinational circuit or synchronous sequential circuits, the table used by the product terms method (MW; PW; FW). Initially, there is a validation procedure, which validates the input variables, the outputs, and the logical

equations. If all syntactic conditions are verified, the expert system generates for each product term a list containing the mask word, the product word, and the function word. Next, the system transforms the set of lists into a single memory list and eliminates the repetition of the words that characterize the same product term. This final list expresses the behaviour of the circuit.

The step 86 uses the tables generated in the step 85 to configure or to re-configure the circuit structure.

The invention can be implemented in digital electronic circuitry, or in computer hardware, firmware, software, or in combinations thereof. Apparatus of the invention can be implemented in a computer program product tangibly embodied in a machine-readable storage device for execution by a programmable processor; and a programmable processor executing a program of instructions by operating on input data and generating output can perform the actions.

The invention can be implemented advantageously in one or more computer programs that are executable on a programmable system including at least one programmable processor coupled to receive data and instructions from, and to transmit data and instructions to, a data storage system, at least one input device, and at least one output device. Each program can be implemented in a high-level procedural or object oriented programming language, or in assembly or machine language if desired; and in any case, the language can be a compiled or interpreted language.

Suitable processors include, by way of example, both general and special purpose microprocessors. Generally, a processor will receive instructions and data from a read-only memory and/or a random access memory.

Generally, the system will include one or more mass storage devices suitable for tangibly embodying computer program instructions and data include all forms of non-volatile memory, including by way of example semiconductor memory devices, such as EPROM, EEPROM, and flash memory devices; magnetic disks such as internal hard disks and removable disks; magneto-optical

disks; and CD_ROM disks. Any of the forgoing can be supplemented by, or incorporated in, ASICs (application-specific integrated circuits).

A method, named the product terms method that allows to implement and/or to change dynamically the logical behaviour of any combinational or synchronous sequential circuits has been presented. The method uses for every product term of logical equations, expressed as a sum-of-product, three memory words: mask word, product word and function word. The words of all product terms are ranged in a table, which characterize the logical behaviour of the circuit.

The invention provides the hardware structure of several new types of VSLI circuits, having re-configurable logic behaviours. A first embodiment implements any type of multiple output combinational circuit, a second embodiment implements any synchronous sequential circuit with only clock input and, a third embodiment implements any synchronous sequential circuits with data inputs and clock input.

An expert system capable to generate the tables used for the product terms method by interpreting and analysing the logical equations either supplied by the user or found in a database is also provided.

Numerous modifications, variations, and adaptations may be made to the particular embodiments of the invention without departing from the scope of the invention which is defined in the claims.

CLAIMS:

1 A VLSI circuit for implementing in hardware any multiple output combinational circuit having the output functions expressed in logical sum-of-product equations, with m inputs, r outputs and n product terms $p(k)$, comprising:

a register with m bits for storing the input variables;

n cells, a cell $C(k)$ for determining the logical value of a product term $p(k)$ of said equations for given inputs;

a logical summing circuit, realised with r OR gates each one with n inputs, associated with said cell $C(k)$ for receiving the logical value of product terms $p(k)$ and outputting the r bits of output functions.

2 A VLSI circuit as defined in claim 1, wherein said cell $C(k)$ comprises:

a storage area for storing the information that characterize a product term, named mask word, product word and function word;

first AND gates means for receiving said inputs and said mask word to produce a first intermediate result, which identify the input variables that form a product term;

EQUIVALENCE gates means for comparing the said product term with said first intermediate result to produce a second intermediate result concerning a product term;

second AND gate means for receiving the said second intermediate result to produce a logical value which is the value of the product term; if this value is logical 1 the product term is active;

third AND gates means to allow passing said function word if said product term is active, and consequently to be OR-ed with function words of other active product terms.

3 A cell $C(k)$ as defined in claim 2, wherein said storage area of a cell $C(k)$ comprises three m -bit registers for m input bits of said $C(k)$ cell.

4 A cell C(k) as defined in claim 2, wherein said first AND gates of a cell C(k) comprises m 2-bit AND gates to produce said first intermediate result for m input bits.

5 A cell C(k) as defined in claim 2, wherein said EQUIVALENCE gates of a cell C(k) comprises m 2-bit EQUIVALENCE ($\overline{\text{XOR}}$) gates to produce said second intermediate result for m input bits.

6 A cell C(k) as defined in claim 2, wherein said second AND gate of a cell C(k) comprises one m-bit AND gate to produce a logical value which is the value of the product term for m input bits.

7 A cell C(k) as defined in claim 2, wherein said third AND gates of a cell C(k) comprises m 2-bit AND gates to allow passing said function word if said product term has the logical value of 1, for m input bits.

8 A VLSI circuit according to claim 2 for implementing in hardware any synchronous sequential circuit with clock input only and outputs taken from the state register, having the next state functions expressed in logical sum-of-product equations, with m bits in state register and n product terms p(k), further comprising:

a clock input;

a state register with m bits for storing the state variables;

n cells of said cell C(k) for determining the logical value of a product term p(k) of said next state equations;

a logical summing circuit, realised with m OR gates each one with n inputs, associated with said cell C(k) for receiving the logical value of product terms p(k) and outputting the m bits of said next state functions;

a feedback connection to establish the next state.

9 A VLSI circuit according to claim 2 for implementing in hardware any synchronous sequential circuit with data inputs and clock input, having the next state functions and the output functions expressed in logical sum-of-product equations, with m inputs, r outputs, s bits in state register, n1 product terms p(k) in next state equations and n2 product terms p(k) in output equations, further comprising:

a clock input;
 a register with m bits for storing the input variables;
 a state register with s bits for storing the state variables;
 $n1$ cells of said cell $C(k)$ for determining the logical value of a product term $p(k)$ of said next state equations;
 $n2$ cells, said cell $C(k)$ for determining the logical value of a product term $p(k)$ of said output equations;
 a logical summing circuit, realised with $(m + s)$ OR gates each one with $n1$ inputs, associated with said cell $C(k)$ for receiving the logical value of product terms $p(k)$ of the next state equations and outputting the $(m + s)$ bits of next state functions.
 a logical summing circuit, realised with r OR gates each one with $n2$ inputs, associated with said cell $C(k)$ for receiving the logical value of product terms $p(k)$ and outputting the r bits of output functions.
 a feedback connection to establish the next state;

10 An expert system for interpreting and analysing the logical behaviour of digital circuits, given by input variables, outputs and logical equations in form of sum-of –products, where interpreting and analysing:

- means to validate said input variables and said outputs;
- means to analyse said logical equations;
- means to generate for every product term said mask word, said product word and said function word;
- means to transform the individual lists, generated for each product term in a single memory list defining the logical behaviour of a digital circuit and storing same in a database.

11 A method for dynamically configuring the logical behaviour of a VLSI circuit as described in claim 1, or a VLSI circuit as described in claim 8, or a VLSI circuit as described in claim 9, by performing the steps of:

inputting the input variables, the outputs and a plurality of sum-of –product equations, which describe the logical behaviour of a digital circuit, to an expert system as defined in claim 10;

generating three memory words uniquely defining each product term of said equations, as defined in claim 10; and

storing said mask word, said product word, and said function word into a corresponding register associated with a cell $C(k)$, as described in claim 2.

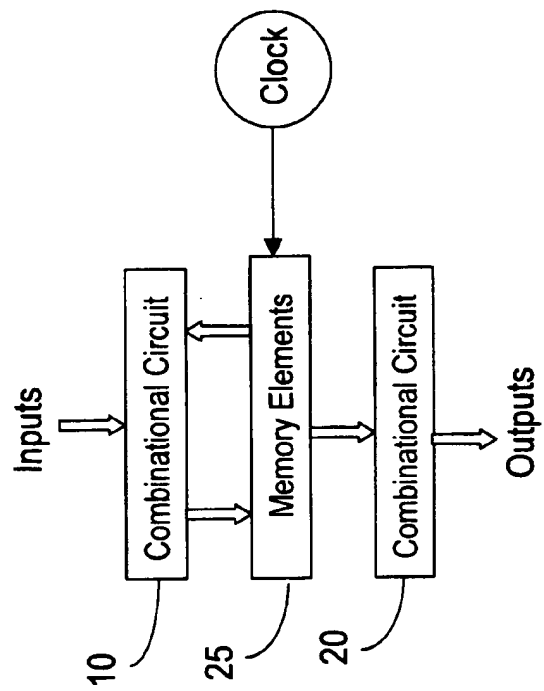


FIGURE 1
(Prior Art)

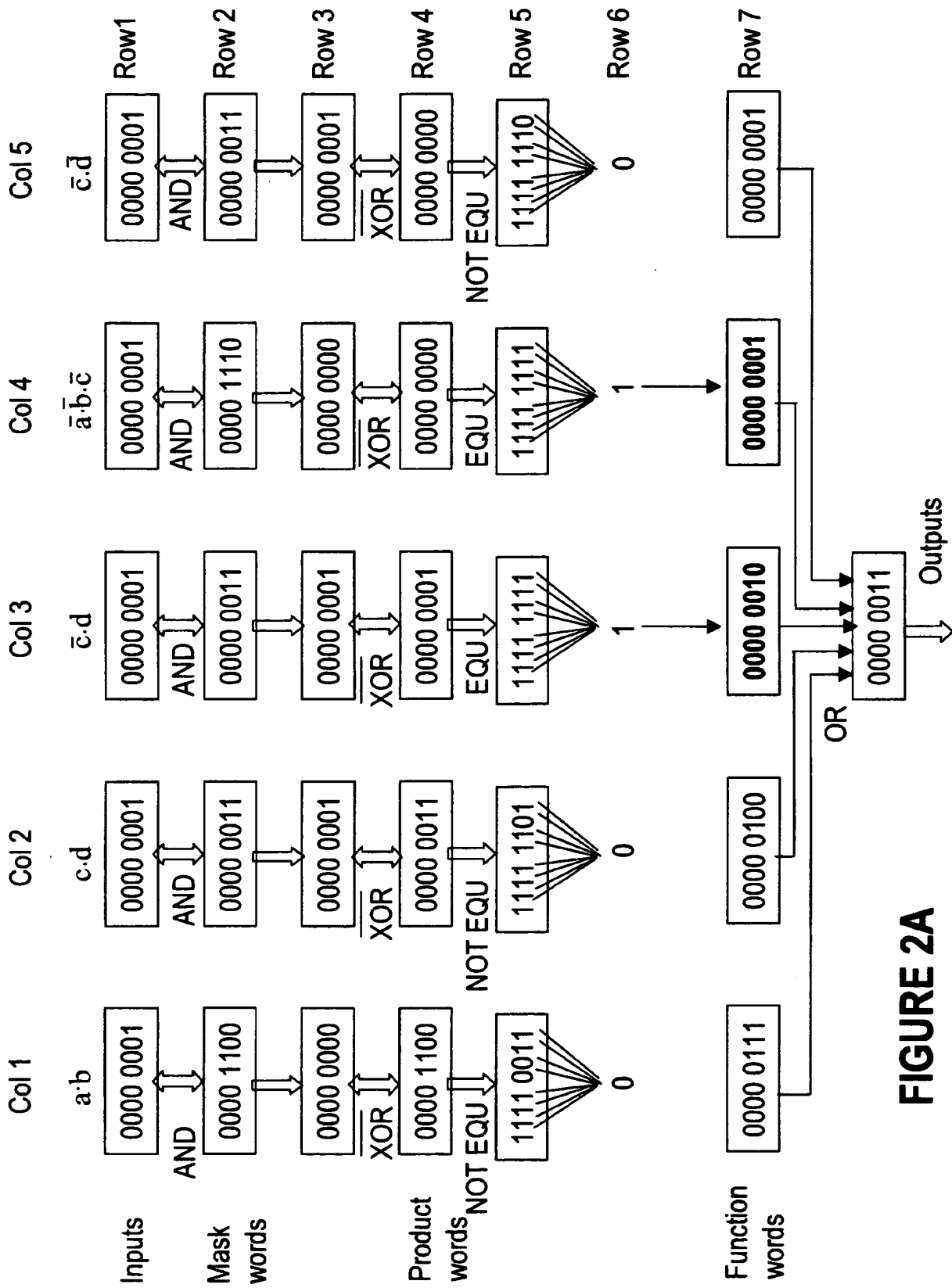


FIGURE 2A

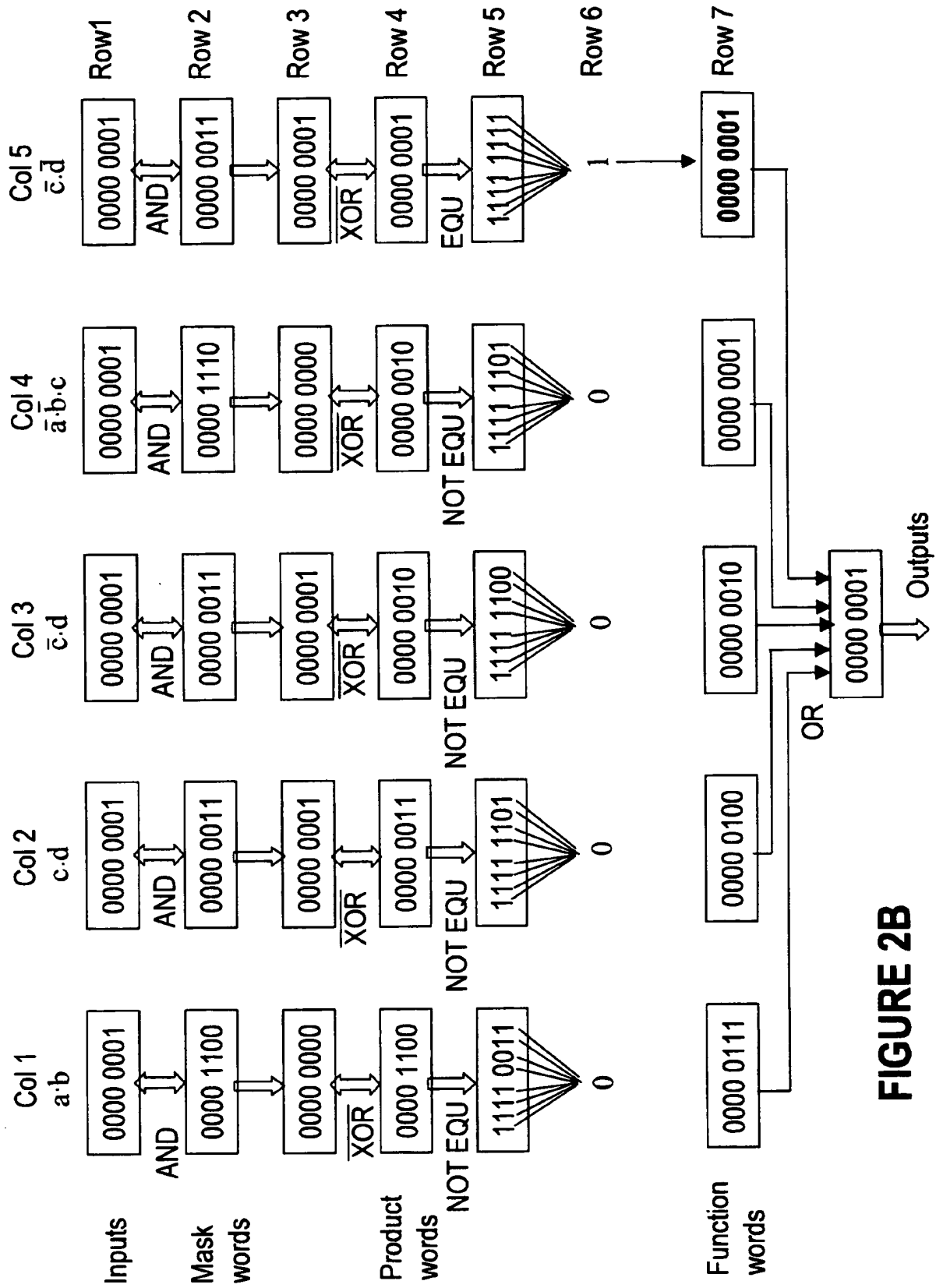


FIGURE 2B

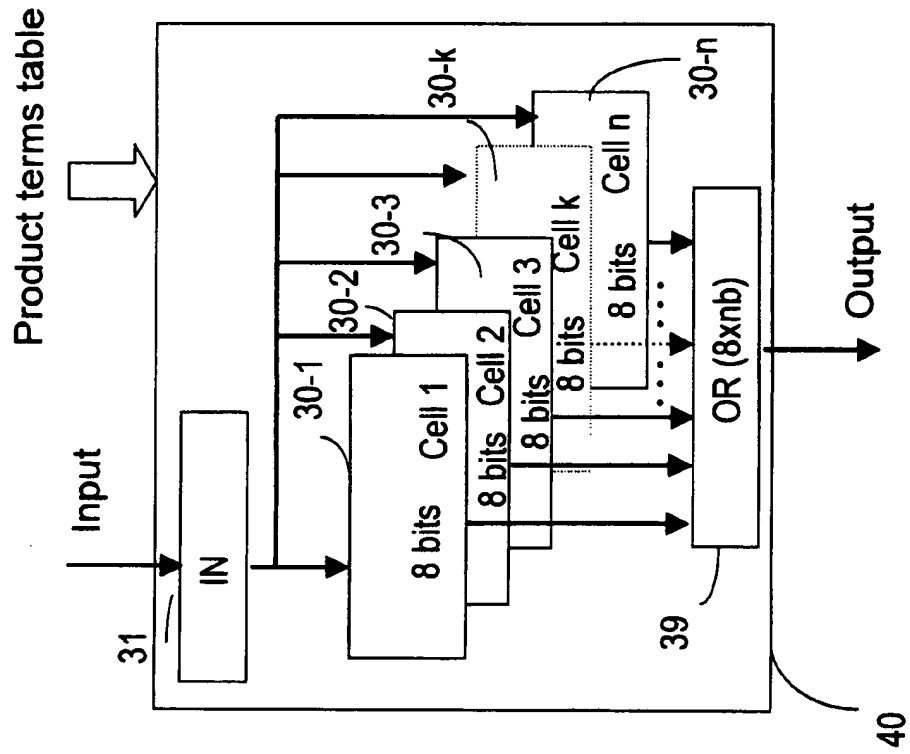
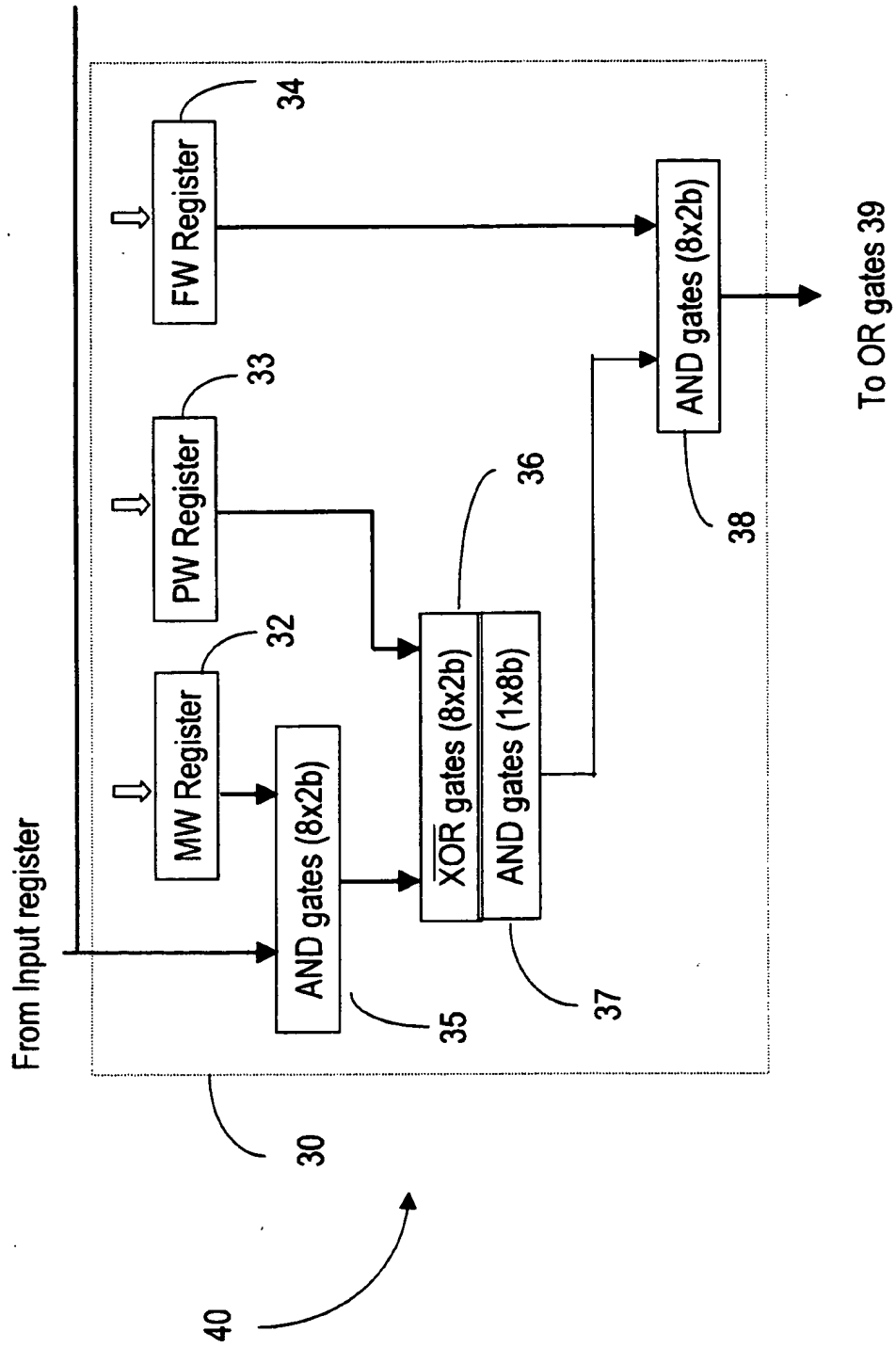


FIGURE 3

**FIGURE 4**

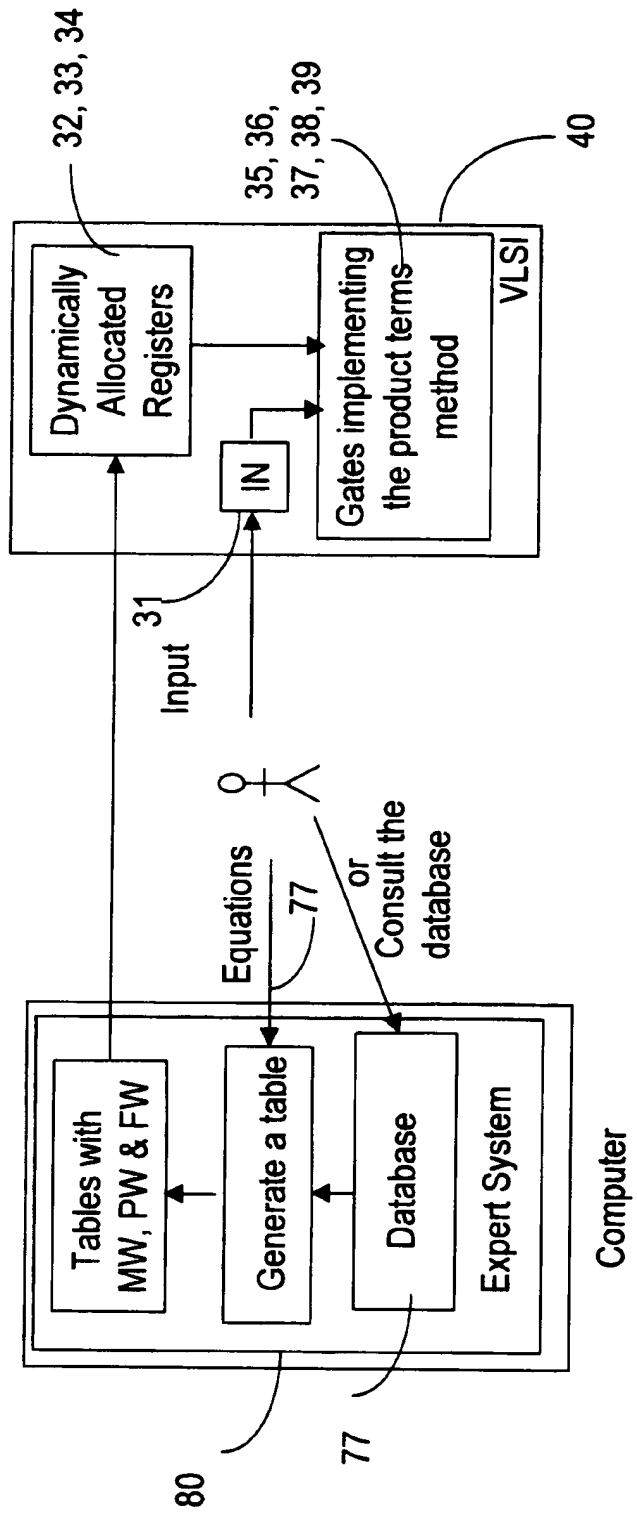
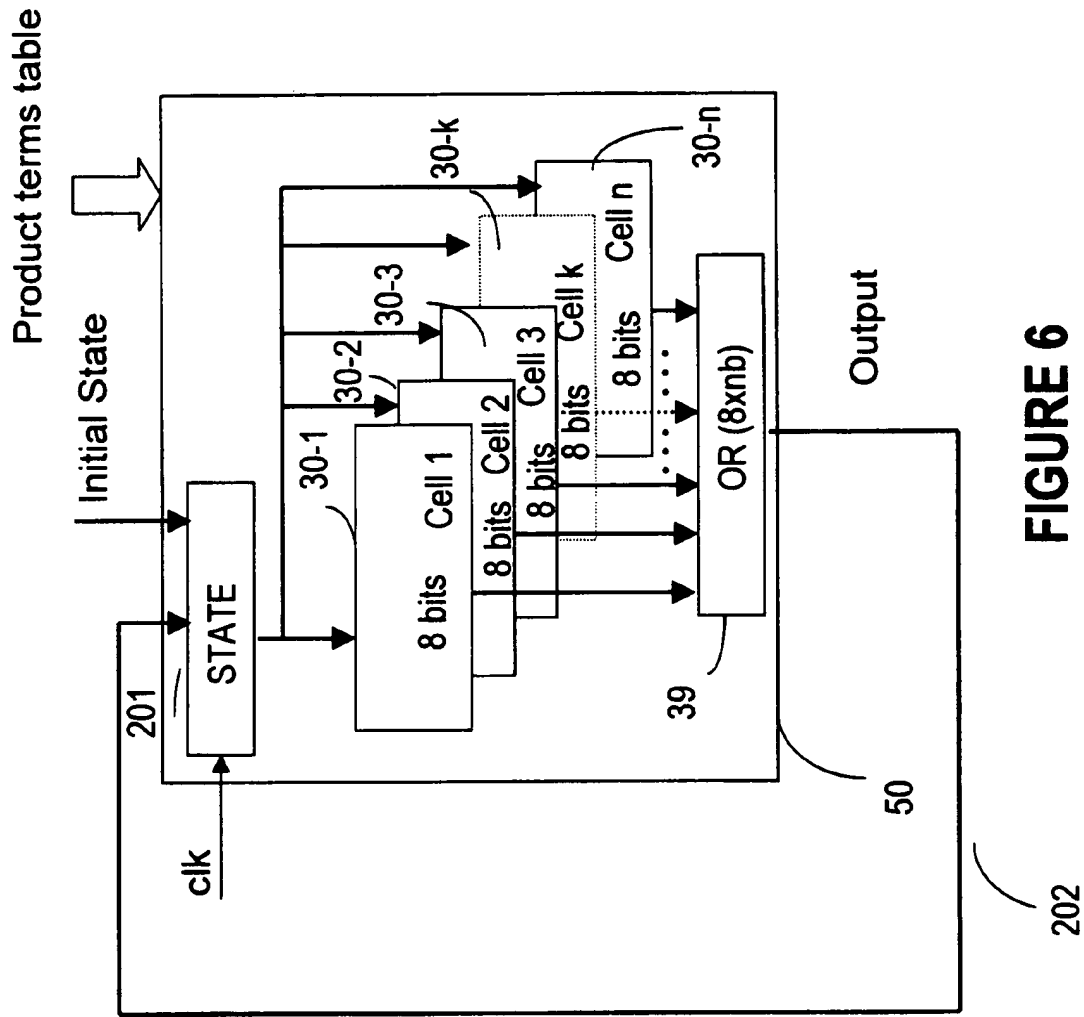
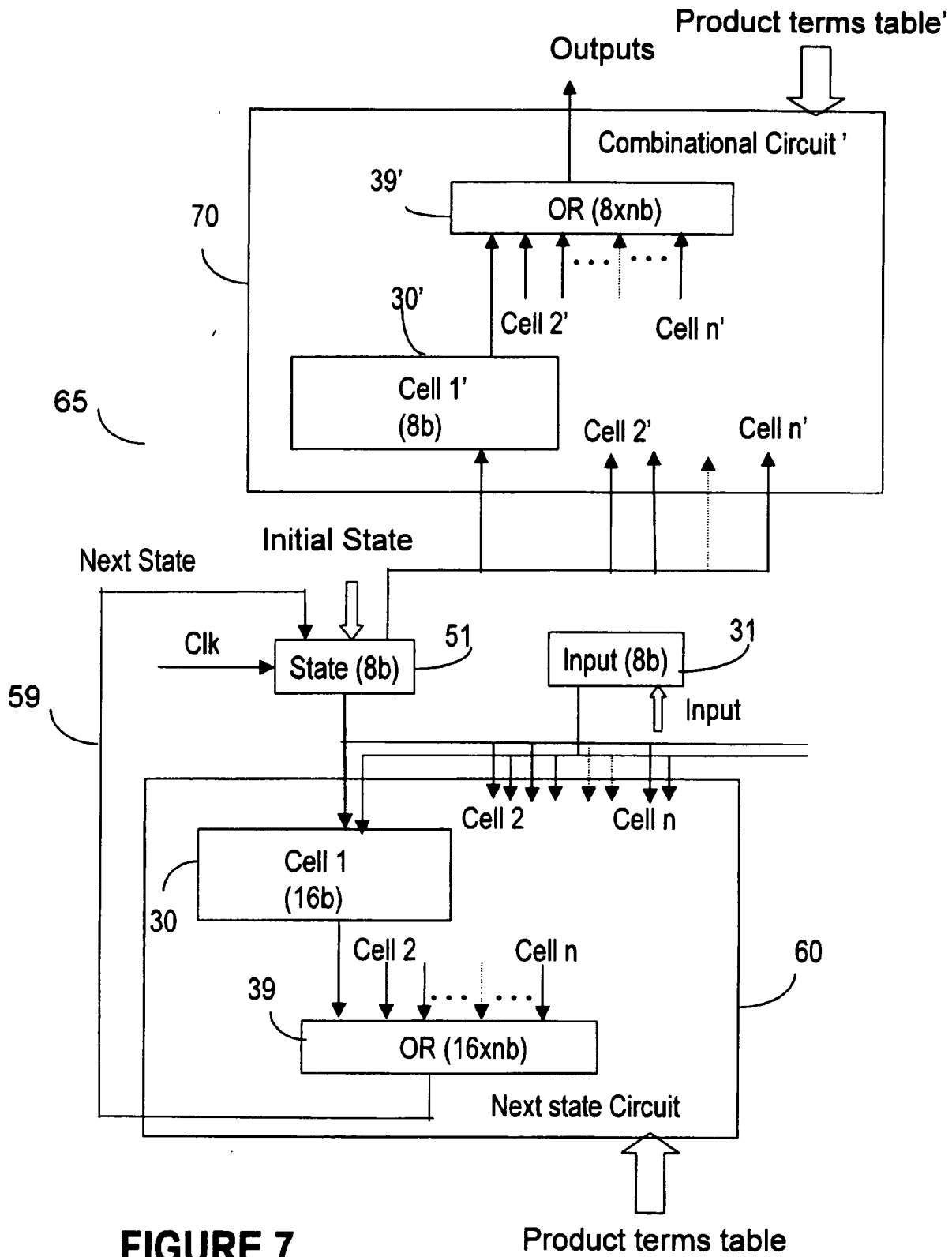
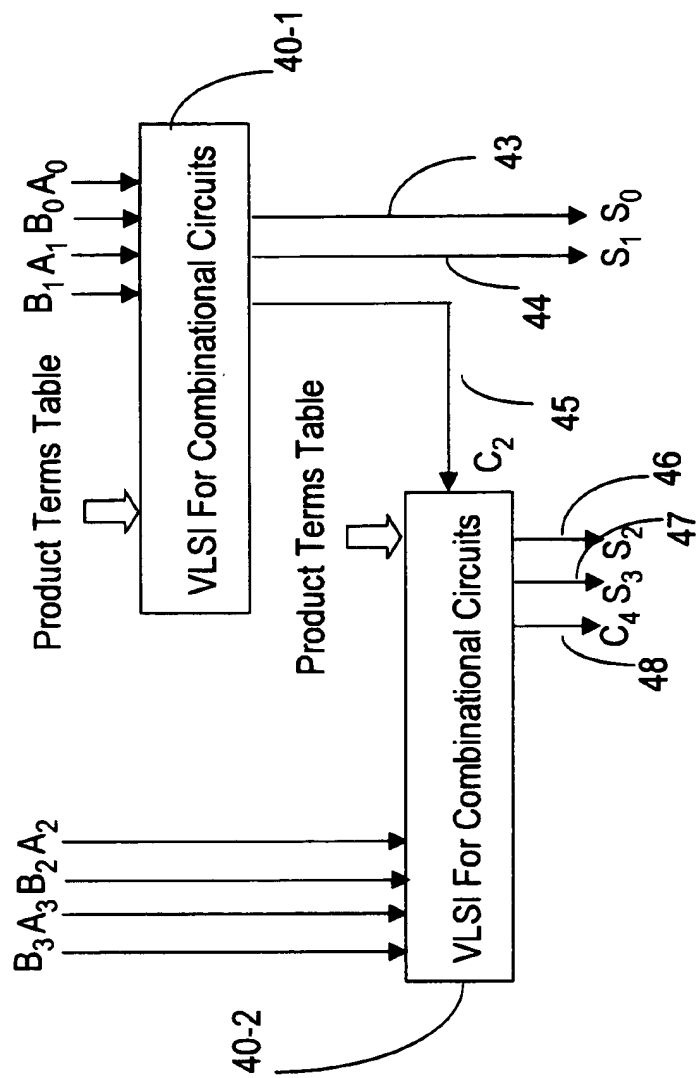


FIGURE 5





**FIGURE 8**

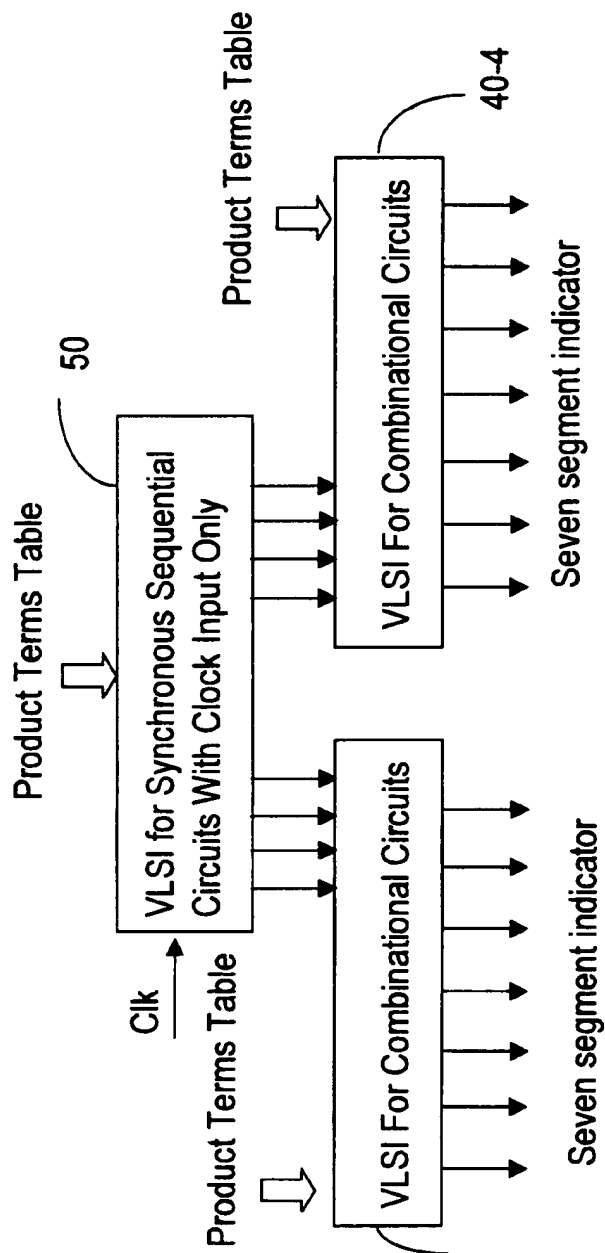
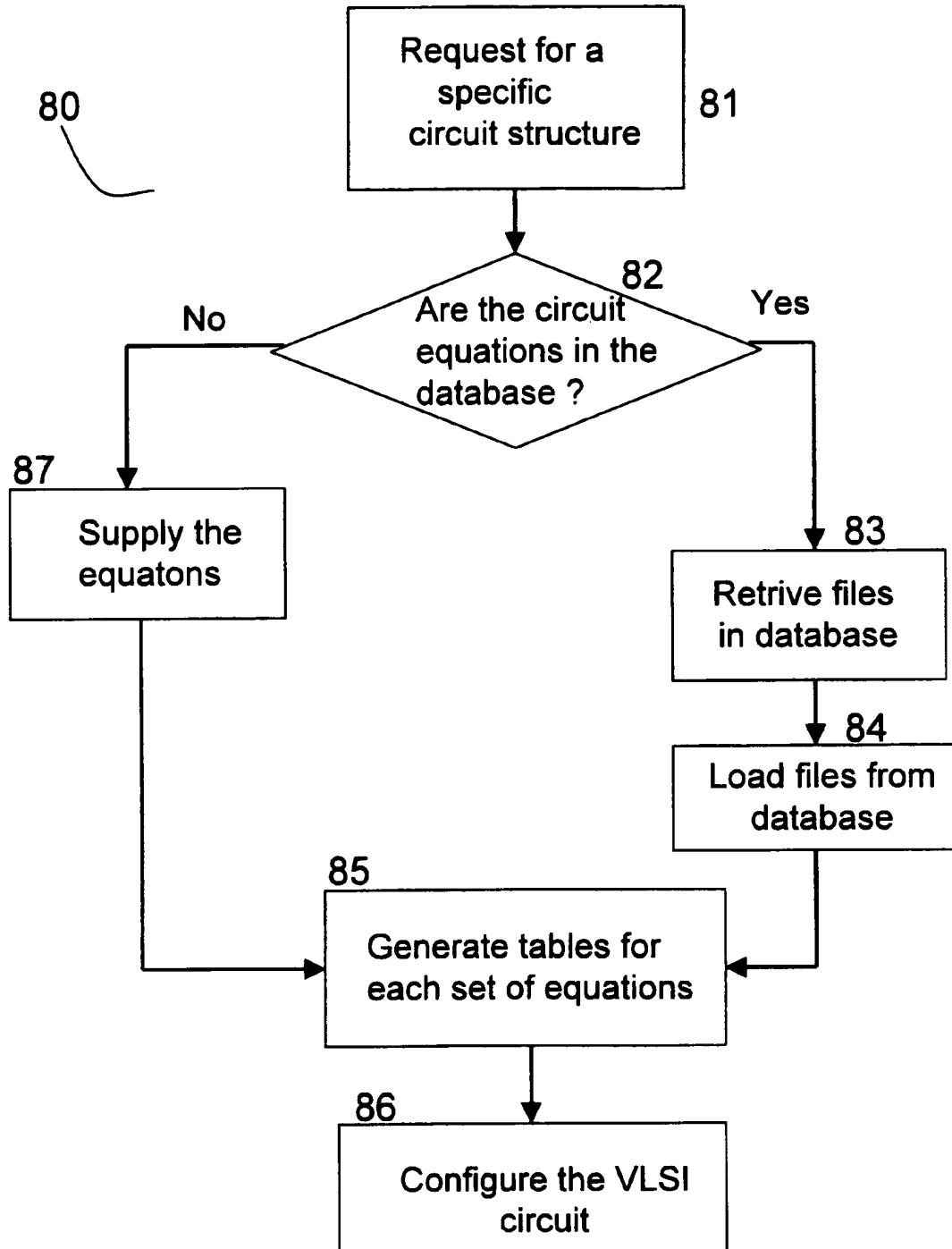


FIGURE 9

**FIGURE 10**